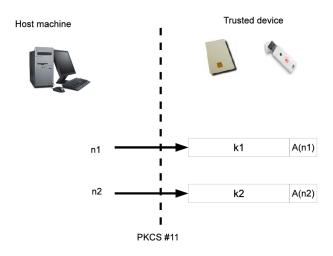
Attacking and fixing PKCS#11

Security Course, Ca' Foscari, 2016

Security APIs



PKCS#11 API for trusted devices

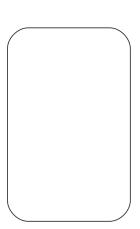


Outline

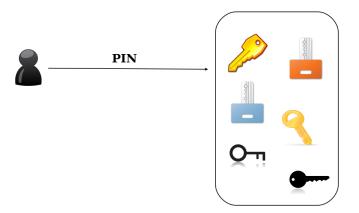
- API Overview
- ▶ PKCS#11 key management attacks
- ► API configuration problems
- ► How to make PKCS#11 secure?

PKCS#11, an overview

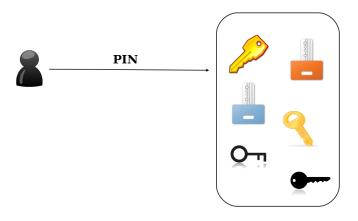




PKCS#11, an overview

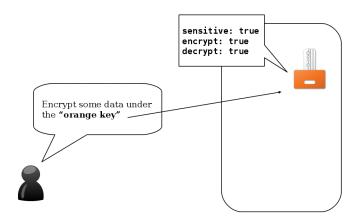


PKCS#11, an overview



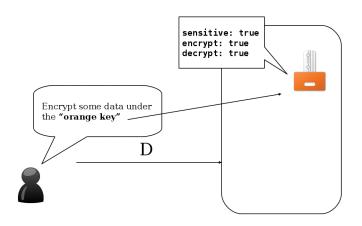
the PIN is a 'second-layer' protection to unlock the token
 ⇒ it should never give access to sensitive key values

PKCS#11 keys and cryptographic operations



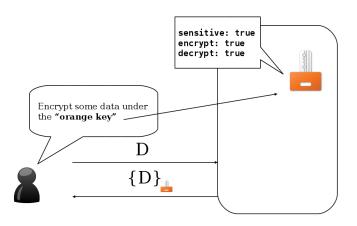
- Keys have attributes and are referenced via handles
- ► APIs for cryptographic operations

PKCS#11 keys and cryptographic operations



- Keys have attributes and are referenced via handles
- ► APIs for cryptographic operations

PKCS#11 keys and cryptographic operations



- Keys have attributes and are referenced via handles
- ► APIs for cryptographic operations

Security of keys

Confidentiality of sensitive keys

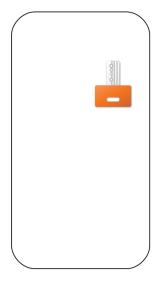
sensitive keys never accessible as plaintext outside the device ... even if we know the PIN

Attack scenario

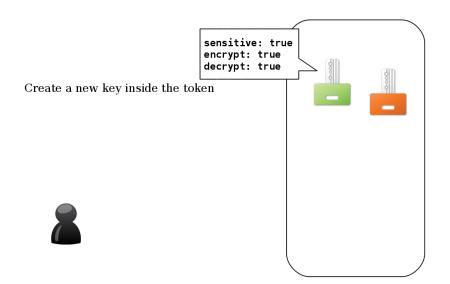
- 1. token used on compromised host
- attacker sniffs PIN and extracts sensitive keys
- 3. attacker clones the token

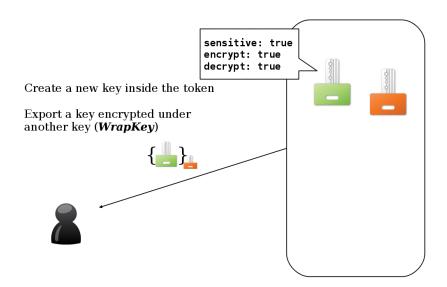
"... the PIN may be passed through the operating system. This can make it easy for a rogue application on the operating system to obtain the PIN ... " [RSA Security]

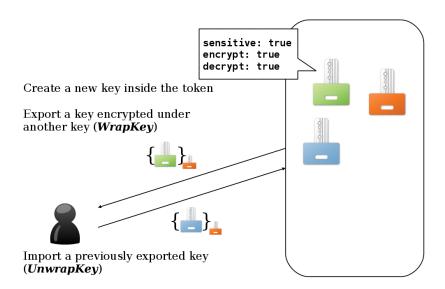
Create a new key inside the token

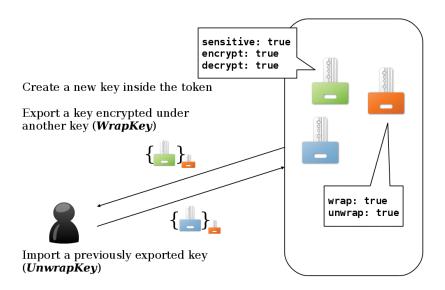




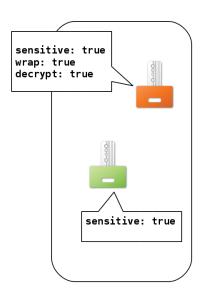


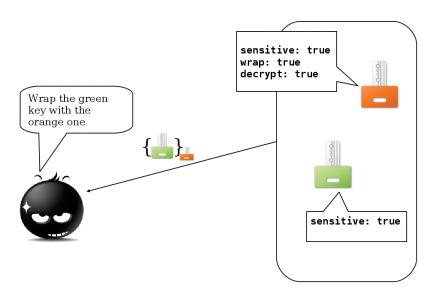


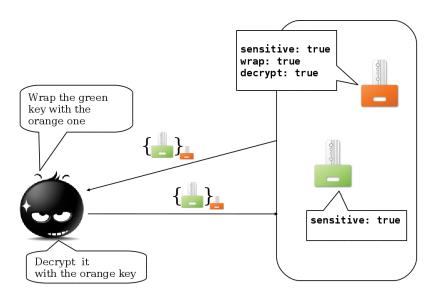


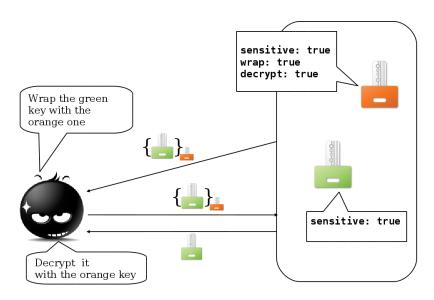




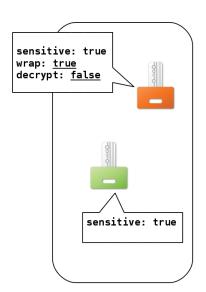


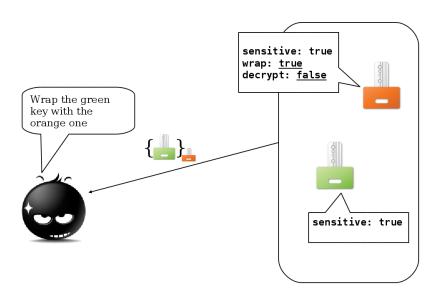


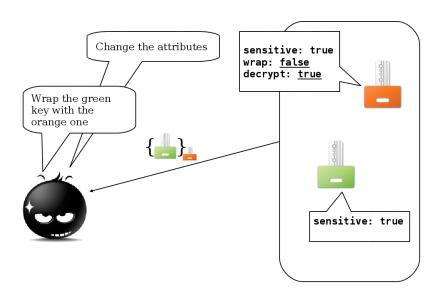


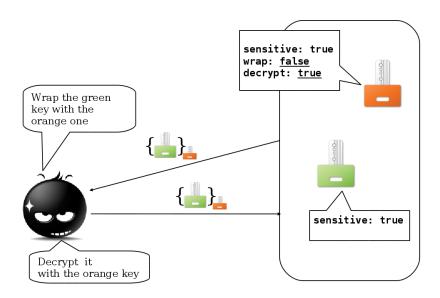


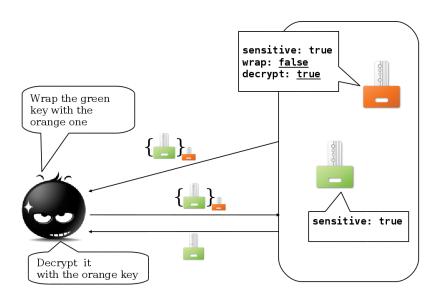






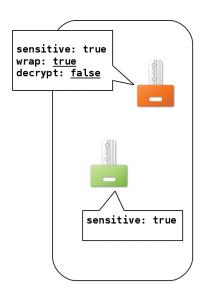




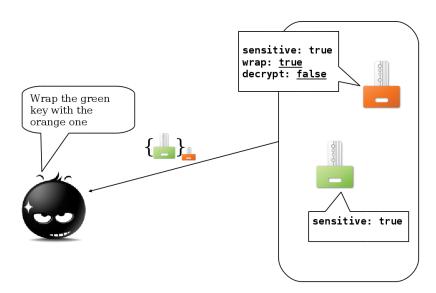


Well ... make attributes 'sticky on'

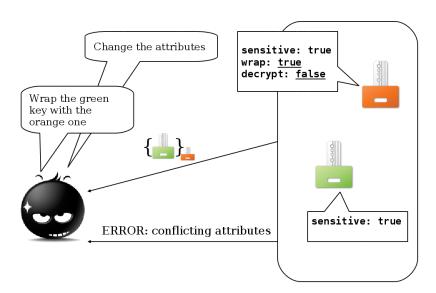


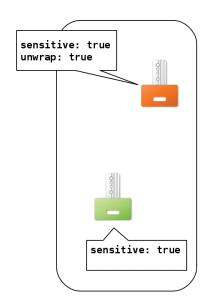


Well ... make attributes 'sticky on'

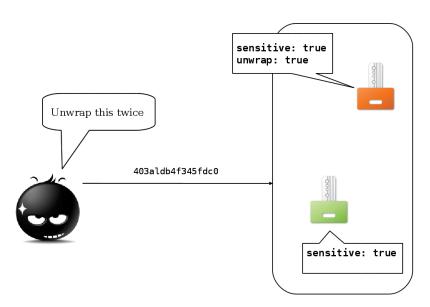


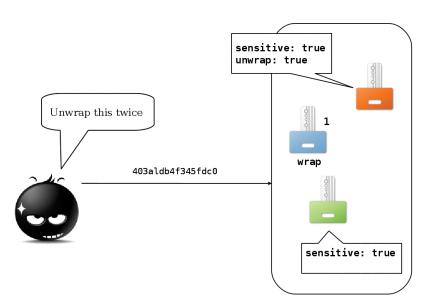
Well ... make attributes 'sticky on'

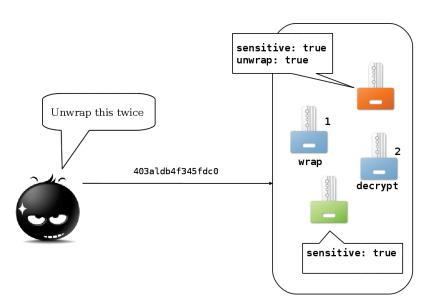


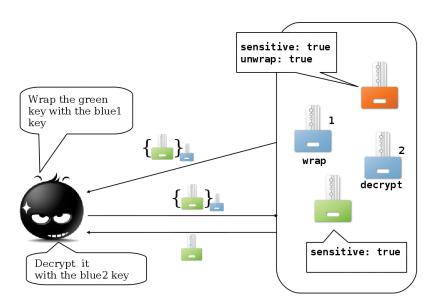












Now what?

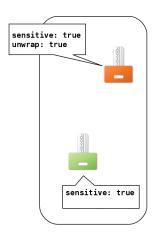
Now what?

- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?

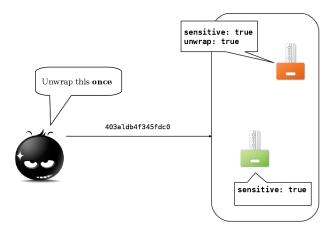
Now what?

- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?

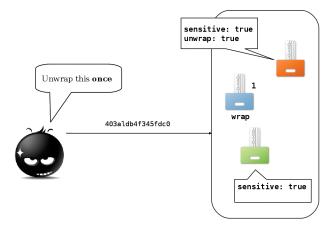




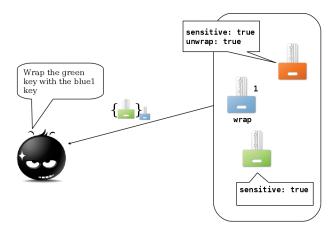
- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?



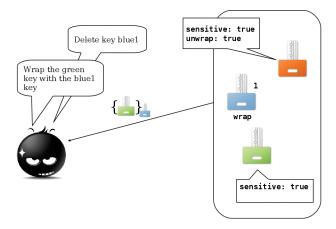
- 😚 check if two instances of the same key have different attributes
 - ▶ is this of any help?



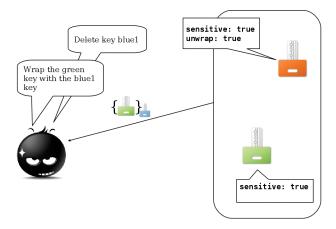
- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?



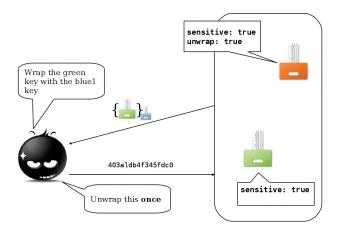
- 😚 check if two instances of the same key have different attributes
 - ▶ is this of any help?



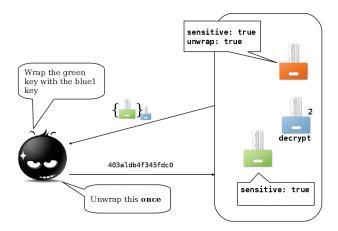
- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?



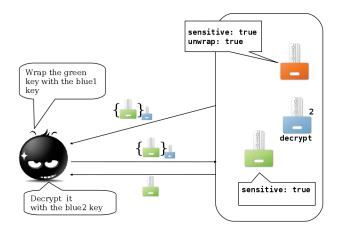
- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?



- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?



- 😵 check if two instances of the same key have different attributes
 - ▶ is this of any help?



Wrapping format

- keep track of key template when wrapping it
- check that it corresponds when unwrapping

Wrapping format

- keep track of key template when wrapping it
- check that it corresponds when unwrapping
- Compute a CBC-MAC of the wrapped key together with its relevant attributes

$$MAC_{k_m}(\{k_1\}_{k_2}, sensitive, wrap, unwrap, ...)$$

and give it as output together with $\{k_1\}_{k_2}$

if the MAC does not correspond the key is not imported

Wrapping format

- keep track of key template when wrapping it
- check that it corresponds when unwrapping
- Compute a CBC-MAC of the wrapped key together with its relevant attributes

$$MAC_{k_m}(\{k_1\}_{k_2}, sensitive, wrap, unwrap, ...)$$

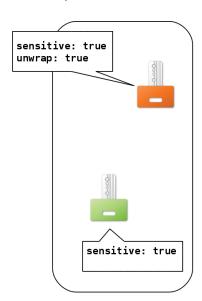
and give it as output together with $\{k_1\}_{k_2}$

▶ if the MAC does not correspond the key is not imported

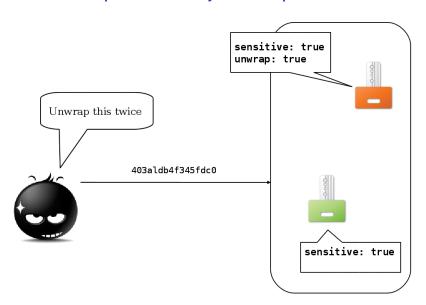
Note: k_m can be derived from k_2 , e.g., by encrypting some constant

Unwrap of arbitrary data is prevented

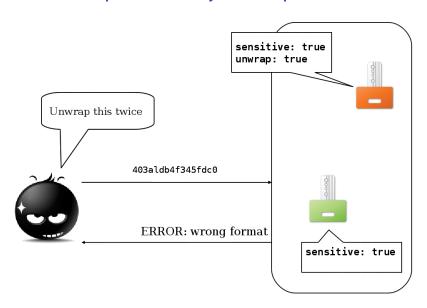




Unwrap of arbitrary data is prevented



Unwrap of arbitrary data is prevented



Summary: Attribute policies and wrapping formats

Sticky

Once an attribute is set (unset), it may not be unset (set). Read-only attributes can be thought as both sticky on and off.

Conflicting

Pairs of attributes that cannot be simultaneously set. (not in the PKCS#11 documentation)

Wrapping format

Keep track of relevant attributes when wrapping, and check they are the same when unwrapping

buffalo buffalo buffalo buffalo buffalo buffalo buffalo

buffalo buffal

buffalo buffalo

buffalo FROM Buffalo WHO buffalo (intimidate) buffalo FROM Buffalo (intimidate) buffalo FROM Buffalo

Formal analysis of PKCS#11 (2 slides!)

► Terms representing keys, ciphertexts, handles

$$k$$
, senc (d, k) , $h(n, k)$

▶ Rules $T; L \xrightarrow{\text{new } \hat{n}} T'; L'$ representing API calls

$$h\left(x_{1},y_{1}\right),y_{2};\;\mathsf{encrypt}\left(x_{1}\right)\;\to\;\mathsf{senc}\left(y_{2},y_{1}\right)$$

▶ Transitions $(S, V) \rightsquigarrow (S', V')$ representing API invocation

```
\langle \; \{h(n,k),d\}; \mathtt{encrypt}(n) \; \rangle \leadsto \langle \; \{h(n,k),d, \frac{\mathsf{senc}(d,k)}{\mathsf{sencrypt}(n)} \; \rangle
```

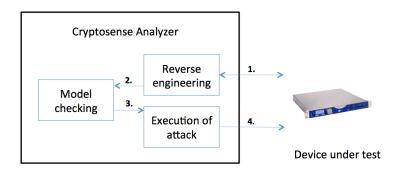
Wrap-Decrypt attack, formally

Rules for key generation, wrap, decrypt:

```
\begin{array}{ccc} & \xrightarrow{\mathrm{new}\;\mathsf{n},\mathsf{k}} & \mathsf{h}\left(\mathsf{n},\mathsf{k}\right); \mathcal{A} \\ \mathsf{h}\left(\mathsf{x}_{1},\mathsf{y}_{1}\right), \mathsf{h}\left(\mathsf{x}_{2},\mathsf{y}_{2}\right); \mathtt{wrap}\left(\mathsf{x}_{1}\right), \mathtt{extract}\left(\mathsf{x}_{2}\right) & \longrightarrow & \mathsf{senc}\left(\mathsf{y}_{2},\mathsf{y}_{1}\right) \\ & \mathsf{h}\left(\mathsf{x}_{1},\mathsf{y}_{1}\right), \mathsf{senc}\left(\mathsf{y}_{2},\mathsf{y}_{1}\right); \mathtt{decrypt}\left(\mathsf{x}_{1}\right) & \longrightarrow & \mathsf{y}_{2} \end{array}
```

- ▶ We start from state $\langle \{h(n_1, k_1)\}, \text{sensitive}(n_1), \text{extract}(n_1) \rangle$ $\rightarrow \langle \{h(n_1, k_1), h(n_2, k_2)\},$
 - sensitive(n_1), extract(n_1), wrap(n_2), decrypt(n_2) $\rangle \sim \langle \{h(n_1, k_1), h(n_2, k_2), senc(k_1, k_2)\}, \langle \{h(n_1, k_1), h(n_2, k_2), senc(k_1, k_2), senc(k_1, k_2)\}, \langle \{h(n_1, k_1), h(n_2, k_2), senc(k_1, k_2), senc(k_1, k_2)\}, \langle \{h(n_1, k_1), h(n_2, k_2), senc(k_1, k_2), senc(k_1, k_2), senc(k_1, k_2)\}, \langle \{h(n_1, k_1), h(n_2, k_2), senc(k_1, k_2$
 - $\langle (n(n_1, n_1), n(n_2, n_2), \text{sens}(n_1, n_2)), \text{sensitive}(n_1), \text{extract}(n_1), \text{wrap}(n_2), \text{decrypt}(n_2) \rangle$
 - $\wedge \langle \{h(n_1, k_1), h(n_2, k_2), senc(k_1, k_2), \frac{k_1}{k_1} \}, \\
 sensitive(n_1), extract(n_1), wrap(n_2), decrypt(n_2) \rangle$

Crytpsense Analyzer



Results of testing of Tokens/Smarcards

Device			Supported Functionality						Attacks found			
Brand	Model	S	as	cobj	chan	W	WS	wd	rs	ru	su	Tk
Aladdin	eToken PRO	√	√	√	√	√	√	√				wd
Athena	ASEKey	✓	\checkmark	\checkmark								
Bull	Trustway RCI	✓	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	✓				wd
Eutron	Crypto Id. ITSEC		\checkmark	\checkmark								
Feitian	StorePass2000	✓	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark		rs
Feitian	ePass2000	✓	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark		rs
Feitian	ePass3003Auto	✓	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark		rs
Gemalto	SEG		\checkmark		\checkmark							
MXI	Stealth MXP Bio	✓	\checkmark		\checkmark							
RSA	SecurID 800	✓	\checkmark	\checkmark	\checkmark				\checkmark	\checkmark	\checkmark	rs
SafeNet	iKey 2032	✓	\checkmark	\checkmark		\checkmark						
Sata	DKey	✓	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark	\checkmark	rs
ACS	ACOS5	√	√	✓	√							
Athena	ASE Smartcard	✓	\checkmark	\checkmark								
Gemalto	Cyberflex V2	✓	\checkmark	\checkmark		\checkmark	\checkmark	✓				wd
Gemalto	SafeSite V1		\checkmark		\checkmark							
Gemalto	SafeSite V2	√	\checkmark	\checkmark	\checkmark	\checkmark	\checkmark	✓	\checkmark	\checkmark	\checkmark	rs
Siemens	CardOS V4.3 B	✓	\checkmark	✓		\checkmark				✓		ru

Towards a secure PKCS#11

Fixing this kind of attacks is far from being trivial

- ► most real tokens are either vulnerable or cut-down so to avoid wrap/unwrap [Bortolozzo et al. CCS'10].
- secure configuration + wrapping format [Delaune et al. JCS'09] (e.g., Eracom mechanism Wrapkey_DES3_CBC, out of the standard)
- more on wrapping formats [Fröschle, Steel ARSPA-WITS'09]
- configuration with mixed roles [Bortolozzo et al. CCS'10]
- configuration with trusted keys [Fröschle, Sommer FAST'11]
- key diversification [Centenaro et al. POST'12]

Summary

- Crypto APIs are irritatingly liberal
- Attacks to compromise a sensitive key and fixes
- ▶ The APIs are hard, sometimes impossible, to configure securely
- Formal methods are effective
 - Find (new) attacks
 - Provide a high level of certification

References



M. Bortolozzo, M. Centenaro, R. Focardi, G. Steel. CryptokiX: a cryptographic software token with security fixes. ASA'10.

M. Centenaro, R. Focardi, F. Luccio.
Type-based Analysis of PKCS#11 Key Management. POST 2012.

J. Clulow. On the security of PKCS#11. CHES'03.

S. Delaune, S. Kremer, G. Steel. Formal analysis of PKCS#11 and proprietary extensions. JCS 2009.

Falcone, A., Focardi R.
Formal Analysis of Key Integrity in PKCS#11. ARSPA-WITS'10.

S. Fröschle and G. Steel. Analysing PKCS#11 Key Management APIs with Unbounded Fresh Data. ARSPA-WITS'09.

S. Fröschle and N. Sommer Concepts and Proofs for Configuring PKCS#11. FAST'11

RSA Security Inc. PKCS #11 v.2.20: Cryptographic Token Interface Standard. June 2004