Denial of Service

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Riccardo Focardi www.unive.it/data/persone/5590470 secgroup.dais.unive.it



Introduction

Denial of Service (DoS)

Increasingly **popular** attack that compromises the **availability** of a service

Example: service is "flooded" by many spurious requests that make it impossible to respond to valid requests

Introduction

Distributed Denial of Service (DDoS)

DoS is particularly effective when launched from **many devices**. Increasing strength in the years ...

- ~400 Mbps in 2002
- ~100 Gbps in 2010
- ~300 Gbps Spamhaus, in 2013
- **~600 Gbps** BBC, in 2015
- ⇒ easily exceed the bandwidth!
 But usually not very long
 (~30min, botnets-for-hire)

Introduction

Distributed Denial of Service (DDoS and IoT)

In 2016 a **new kind of attack** on Dyn, a DNS provider

- long, many hours
- involved multiple attacks from over 100,000 devices
- loT (Internet of Things)
 devices, such as webcams and
 baby monitors
- reached a ~1.2 TBps peak

Definition

NIST SP 800-61

DoS: An attack that **prevents** or **impairs** the **authorized use** of networks, systems, or applications by **exhausting resources**

Targets:

- network bandwidth
- system resources
- application resources

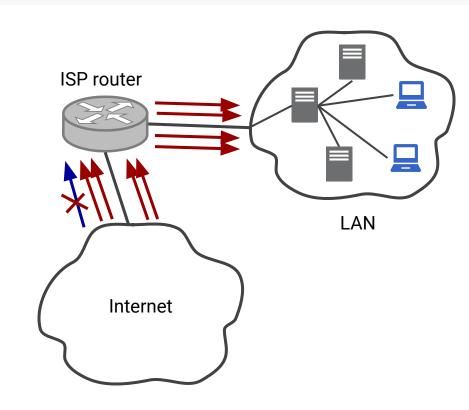
Target: network bandwidth

Network bandwidth: capacity of links connecting a server to the Internet

⇒ Usually the link to the Internet Service Provider (ISP)

If incoming traffic exceeds the bandwidth packets will be discarded

→ Legitimate packets discarded if malicious ones exceeds network bandwidth



Target: system resources

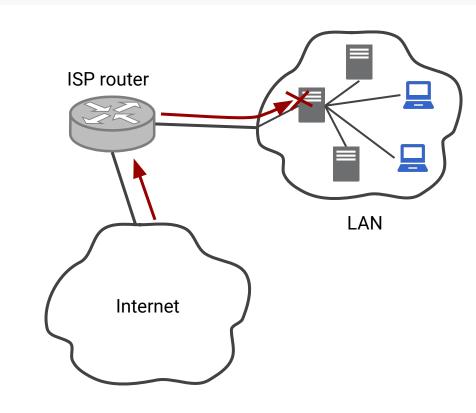
Network handling resources:

required to implement network protocols (e.g. buffers)

⇒ When limit is reached new network connections are refused
Example: TCP connections

"Poison packets" might trigger bugs that break network services

Examples: ping of death, teardrop



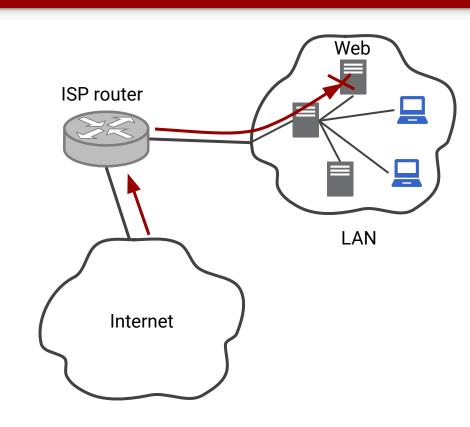
Target: application resources

Application resources: required to accomplish tasks

 ⇒ When limit is reached application becomes unresponsive
 Example: Excessively complex queries to a database

Destructive attacks, exploiting bugs, that crash the application

Examples: Piggybacked SQLi



Target

- Network bandwidth
- System resources
- Application resources

Flooding attacks

Flooding attack: overwhelm network capacity

ICMP flooding: the Internet Control
Message Protocol (ICMP) is used to
send error messages and operational
information

Example: **ping** allows for testing connectivity (-f option **floods the server**)

UDP flooding: attacker targets a UDP service (es. diagnostic echo)

TCP flooding: attacker targets TCP services

Target server might either **respond**, generate a ICMP **destination unreachable** or **reject** the packet

⇒ Responses <u>increase the load!</u>

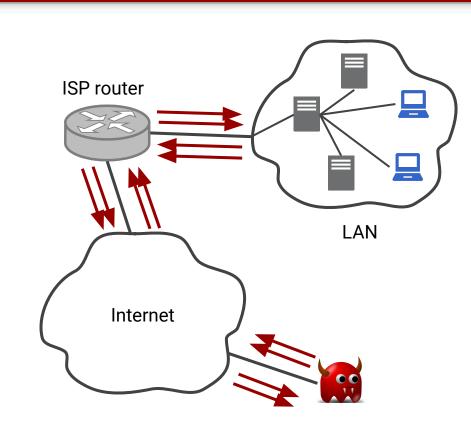
Simple flooding

Simple flooding attack: overwhelm network capacity from a single host

Example: ping with -f option

However:

- Source easily identified (legal actions taken)
- Target will respond "reflecting" the attack back



Source address spoofing

Source address spoofing: attacker use *raw socket interface* to change source address

Randomly selected source addresses

- ⇒ Responses will be scattered around the Internet
- ⇒ Possible errors packets from spoofed address will go towards the target and contribute to DoS

Source address spoofing makes it hard to identify the attacker

Cause: TCP/IP <u>does not ensure</u> that source address really corresponds to the originating host

It would be necessary to (manually) query the logs of **traversed routers** in order to identify the trajectory

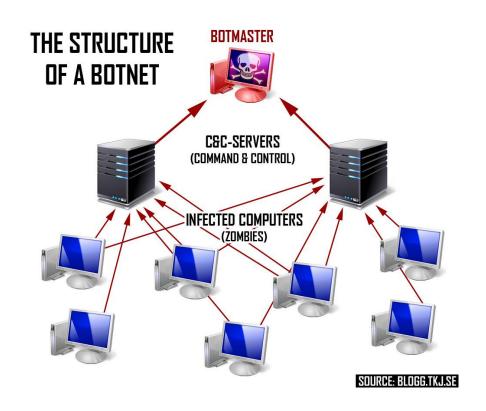
Distributed DoS (DDos)

Botnets: a collection of **zombie devices** under the attacker control

Botnets <u>are hired</u> for DDos

 ~40% of DDoS in 2015 were from botnets for hire

Flooding coming from thousands of hosts easily reaches **Gbps** bandwidth



Target

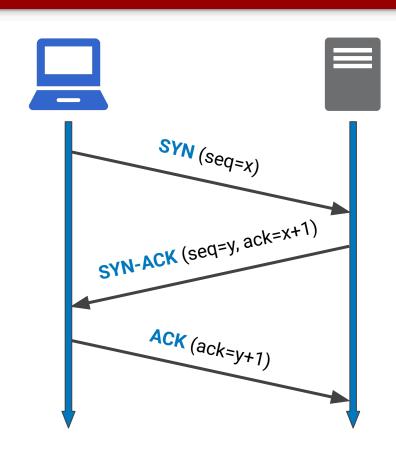
- Network bandwidth
- System resources
- Application resources

SYN spoofing

This attack **overflows** the tables used to manage TCP connections

TCP uses a **three-way handshake** to establish a connection:

- IP lost packets are transparently resent
- Applications using TCP won't notice lost packets and retransmissions



SYN spoofing (ctd)

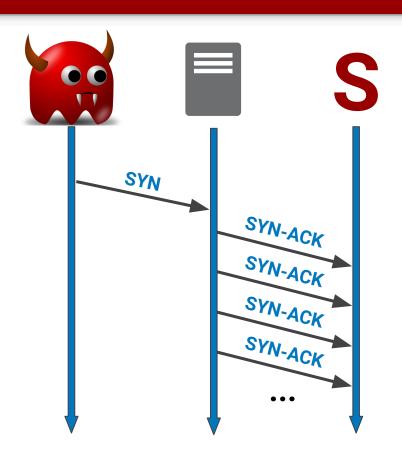
Attack scheme:

- Attacker sends SYN packets with spoofed source addresses
- 2. For each spoofed source S:
 - a. Server sends SYN-ACK to S
 - b. If Server timeouts and N_s <MAX

$$N_s = N_s + 1$$

goto a

- c. Delete connection with S
- ⇒ Table of TCP requests overflows



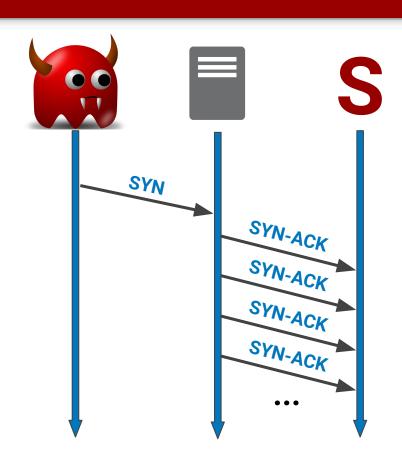
SYN spoofing (ctd)

Attacker sends **enough forged requests** to keep the table full

⇒ Server is **cut off** from the Internet

NOTE: Using **random** spoofed addresses make the probability of not getting a **RST** (reset) answer high

The volume of **SYN** requests is **low** and far from link capacity



Target

- Network bandwidth
- System resources
- Application resources

Application protocol flooding

Attacker floods an application protocol

Examples:

- Session Initiation Protocol (SIP)
 used in VoIP. INVITE requests go
 through proxies and consume
 system/network resources
- HTTP requests can be heavy (e.g. download of large file)

Slowloris: a particular DoS attack that leverages server multi-threading

- start many HTTP requests
 without completing them
- keep the connection alive by sending new lines, periodically

Consumes all available web server connections (in terms of internal system/application resources)

DoS techniques

- Reflection
- Amplification

Reflection attacks

Attacker sends packets to an intermediary with a spoofed source address of the target

The *intermediary* responds to the actual *target*

- If response is larger than request, attack is also amplified
- Tracing is hard if attacker uses many intermediaries

Examples: DNS, SNMP and **ISAKMP** has been exploited for reflection (they can generate **large** responses)

TCP SYN reflection: Attacker can send SYN so that intermediary sends SYN-ACK which in turns generates a RST packet

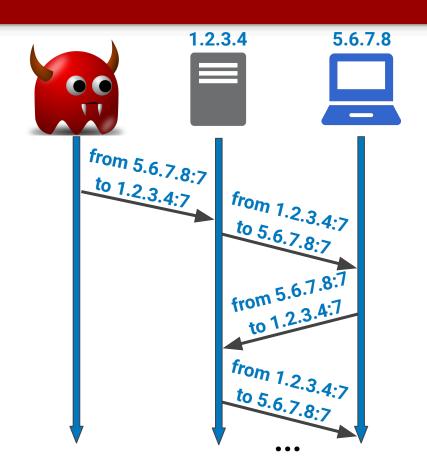
⇒ Both SYN-ACK and RST flood target's network

Reflection "loop"

When echo service (port 7) is enabled reflection loops are possible

Example: The attacker sends a packet to **1.2.3.4** port **7**, with spoofed source address **5.6.7.8** port **7**

- Intermediary echoes to target
- Target echoes to Intermediary
- ... (loop)

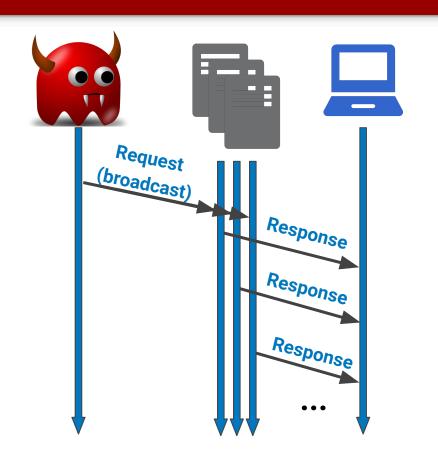


Amplification attacks

Amplification: generating multiple response packets with a single request

Example: send a packet to the **broadcast** address of a network with spoofed address

→ all hosts (with the service enabled)
will respond to the target



Defenses

DoS cannot be fully prevented: attackers that can flood a service with legitimate requests will limit traffic from other users

DoS can be "incidental": important **news** make legitimate users overload referenced web sites

Defenses

what and when

Prevention and mitigation

(**before** the attack)

Detection

(during the attack)

Source traceback

(during and after the attack)

Reaction

(after the attack)

Prevention: spoofed source addresses

Solution 1: **filtering** spoofed source address as close as possible to the originating host

Example: where the organization's network **connects** to the Internet

Filtering spoofed source addresses is a standard security recommendation (RFC 2827) which is too often disregarded!

Solution 2: ensure that the **path back** to the claimed source address is the one being used by the current packet

Example: CISCO implemented this however when **routing is** asymmetrical (path $A \rightarrow B$ and $B \leftarrow A$ differ) this solution is too strict

Multihoming: Connecting to many networks for reliability/performance

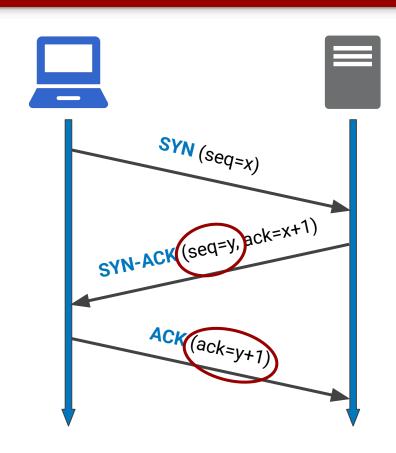
Prevention: SYN spoofing attack

Make the protocol "stateless" by encoding state information directly in the SYN-ACK sequence number y

→ No TCP request table overflow!

When the ACK y+1 is received the server can reconstruct state information from y

Example: "SYN Cookies" in FreeBSD and Linux (similar idea in Windows)



Mitigation: rate limits and random drops

ICMP and UDP flooding to diagnostic services can be mitigated by imposing **limits on packet rates**

Similarly, SYN spoofing attacks can be mitigated by limiting the **connection rate** to a certain service Table overflow of SYN spoofing can be mitigated by **randomly dropping** connections

IDEA: overflow is a probable sign of attack, randomly dropping a connection will more likely **drop an attacker's connection**

Could drop a legitimate connection but it is **better than full DoS**

Other prevention techniques

Block broadcast (amplification)

Block/limit suspicious services and combinations of ports (reflection)

Check **human interaction**, e.g. with captcha (**application** resources)

Keep systems up-to-date and secured (do not become part of a botnet)

Monitor systems, especially high-performance, well-connected servers (potential **intermediary**)

Use **mirrored** and **replicated** servers to increase reliability and **resilience** to DoS attacks

Detection, source traceback and reaction

Detection: capturing packet flows and analyzing them. If the attack is identified

- suitable filters can be activated
- bugs can be fixed
- alternate backup servers can be activated
- ..

Source traceback: necessary for **legal actions**, need collaborating ISPs (can be complex)

Reaction: analyzing the attack and response in order to gain benefit from the experience and to improve future handling. The organization's security can be **improved** as a result